

# Formalization of Foundations of Geometry

## An overview of the GeoCoq library

Julien Narboux

Unité de formation et de recherche  
de **mathématique** et d'**informatique**  
Université de Strasbourg

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## 1 Interactive Theorem Proving for the Education

## 2 Overview of GeoCoq

- Foundations
- Arithmetization of Geometry
  - Addition
  - Multiplication
- Automation
- Continuity
- Logic
- 34 parallel postulates
- Two formalizations of the Elements
- Some high-school examples

# Using a computer to teach maths

## Software are used in the classroom

- 1 for numerical computations
- 2 for symbolic computations (Maple...)
- 3 for producing conjectures (dynamic geometry software)
- 4 for construction exercises (Euclidea, ...)
- 5 for checking conjectures using probabilistic or algebraic methods (Cabri, GeoGebra, ...)

**But**  
the use of software for checking **proofs** is not widespread !

# Teaching the concept of proof

- Maths teachers often do not know about logic.
- Only some of the reasoning rules are given: proof by contradiction, contrapositive, reasoning by cases.
- Semantics checks are used rather than syntactic checks.
- Learning by imitation (a proof is what makes the teacher happy/good marks).

"If you can't explain mathematics to a machine, it is an illusion to think you can explain it to a student."  
De Bruijn "Invited lecture at the Mathematics Knowledge Management Symposium", 25-29 November 2003, Heriot-Watt University, Edinburgh, Scotland



# Different potential goals

- To teach what is a proof
- To teach logic
- To teach software foundations
- To automate proof checking
- To teach maths in general
- To automate feedback in general

# Different potential goals

- To teach what is a proof
- To teach logic
- To teach software foundations
- To automate proof checking
- ~~To teach maths in general~~
- ~~To automate feedback in general~~

I do not need a tutor, just a proof checker.

# ITP, why ?

- Clarify the rules of the game: the deduction rules are explicit.
- Clarify the language: axiom, theorem, lemma, hypotheses, definition, conjecture, counter-example. . .
- Objective criterion for the validity of a proof.
- Interactivity: feedback during homework.
- Motivation: theorem proving as a game.



## Challenges

- Find a good language/user interface.
- Build the needed libraries.
- Automate what should be automatized and not more (depending on the context).

# About the language, proof rules, user interface

I would like deduction rules which are:

- sound
- explicit
- clear
- complete
- not necessarily minimal
- not too far from the mathematical practice

# Coherent logic

$$\forall x, H_1(x) \wedge \dots \wedge H_n(x) \rightarrow \begin{matrix} \exists y, P_1(x, y) \wedge \dots \wedge P_k(x, y) \\ \vee \\ \dots \end{matrix}$$

Several authors have identified independently this fragment of FOL. Allows proofs to be **somewhat** readable <sup>1</sup>.

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<sup>1</sup>Sana Stojanović et al. (2014). “A Vernacular for Coherent Logic”. English. In: [Intelligent Computer Mathematics](#). Vol. 8543. Lecture Notes in Computer Science

# Existing tools

Two communities:

- 1 Didactics of mathematics
- 2 Interactive theorem proving

# Didactics of mathematics Community

Geometry Tutor <sup>2</sup>, MENTONIEZH <sup>3</sup>, DEFI <sup>4</sup>, CHYPRE <sup>5</sup>, Geometrix <sup>6</sup>, Cabri Euclide <sup>7</sup>, Baghera <sup>8</sup>, AgentGeom, geogebraTUTOR and Turing <sup>9</sup>

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<sup>2</sup>John R. Anderson, C. F. Boyle, and Gregg Yost (1985). “The geometry Tutor”. In: [IJCAI Proceedings](#)

<sup>3</sup>Dominique Py (1990). “Reconnaissance de plan pour l’aide à la démonstration dans un tuteur intelligent de la géométrie”. [PhD thesis. Université de Rennes](#)

<sup>4</sup>Ag-Almouloud (1992). “L’ordinateur, outil d’aide à l’apprentissage de la démonstration et de traitement de données didactiques”. [PhD thesis. Université de Rennes](#)

<sup>5</sup>Philippe Bernat (1993). [CHYPRE: Un logiciel d’aide au raisonnement](#). [Tech. rep. 10. IREM](#)

<sup>6</sup>Jacques Gressier (1988). [Geometrix](#).

<sup>7</sup>Vanda Luengo (1997). “Cabri-Euclide: Un micromonde de Preuve intégrant la réfutation”. [PhD thesis. Université Joseph Fourier](#)

<sup>8</sup>Nicolas Balacheff et al. (1999). [Baghera](#).

<sup>9</sup>Philippe R. Richard et al. (2011). “Didactic and theoretical-based perspectives in the experimental development of an intelligent tutorial system for the learning of geometry”. en. In: [ZDM 43.3](#)

# ITP Community

## Computer Science

- logic
- proof of programs, semantics, software foundations

U-Penn, Portland, Princeton, Harvard, Warsaw, CNAM, Lyon, Nice, Paris, Strasbourg, ...

## Maths

- Bachelor - Logic: Bordeaux, Warsaw, Pohang, Strasbourg, ...
- Bachelor - Maths: Nijmegen (ProofWeb), Nice (CoqWeb), ...
- ...

Two kinds of systems:

- 1 Syntactic sugar added over a state of the art proof assistant
  - ▶ PCoq <sup>10</sup>
  - ▶ Coq Web <sup>11</sup>
  - ▶ ProofWeb <sup>12</sup>
  - ▶ Edukera <sup>13</sup>
- 2 Natural language + Automatic Theorem Proving

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<sup>10</sup> Ahmed Amerkad et al. (2001). “Mathematics and Proof Presentation in Pcoq”. In: [Workshop Proof Transformation and Presentation and Proof Complexities in connection v](#)

<sup>11</sup> Jérémy Blanc et al. (2007). “Proofs for freshmen with Coqweb”. In: [PATE'07](#)

<sup>12</sup> CS Kaliszyk et al. (2008). “Deduction using the ProofWeb system”. In:

<sup>13</sup> Benoit Rognier and Guillaume Duhamel (2016). “Présentation de la plateforme edukera”. In:

## Two kinds of systems:

- 1 Syntactic sugar added over a state of the art proof assistant
- 2 Natural language + Automatic Theorem Proving
  - ▶ SAD <sup>10</sup>
  - ▶ Naproche <sup>11</sup>
  - ▶ Lurch <sup>12</sup>
  - ▶ ELFE <sup>13</sup>
  - ▶ CalcCheck <sup>14</sup>
  - ▶ Mendes' system <sup>15</sup>

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<sup>10</sup>[Alexander Lyaletski, Andrey Paskevich, and Konstantin Verchinine \(2006\)](#). “SAD as a mathematical assistant—how should we go from here to there?” In: [Journal of Applied Logic. Towards Computer Aided Mathematics 4.4](#)

<sup>11</sup>[Marcos Cramer et al. \(2010\)](#). “The Naproche Project Controlled Natural Language Proof Checking of Mathematical Texts”. In: [Controlled Natural Language](#)

<sup>12</sup>[Nathan C. Carter and Kenneth G. Monks](#). “Lurch: a word processor built on OpenMath that can check mathematical reasoning”. In:

<sup>13</sup>[Maximilian Doré \(2018\)](#). “The ELFE Prover”. In: [25th Automated Reasoning Workshop](#)

<sup>14</sup>[Wolfram Kahl \(2018\)](#). “CalcCheck: A Proof Checker for Teaching the “Logical Approach to Discrete Math””. In: [Interactive Theorem Proving. Lecture Notes in Computer Science](#)

<sup>15</sup>[Alexandra Mendes and João F. Ferreira \(2018\)](#). “Towards Verified Handwritten   



# Edukera (Rognier and Duhamel)

- Web-application
- Coq is hidden inside the web-page
- LCF style interaction + proof displayed in a pen and paper style.
- Some users in France (about 1000 students, 70k exercises)
- No textual input "proof by pointing", syntactically correct by construction (as using Scratch)
- Easy to learn using a tutorial
- Always correct applications of a logic rule
- Meta-variables

# Two modes

## 1 Logic

- ▶ Use natural deduction rules.
- ▶ Can display proof tree (Fitch's or Gentzen's style).
- ▶ Backward reasoning

## 2 Maths

- ▶ Forward/Backward reasoning.
- ▶ Less fine grained proof steps than in logic mode.

# Edukera (logic mode)



Implication		
$\Rightarrow$	Introduction ( $\Rightarrow I$ ) $\uparrow$	$\mathcal{Q}$
$\Rightarrow_x$	Elimination ( $\Rightarrow E$ ) $\uparrow \downarrow$	$\mathcal{Q}$
Conjunction		
$\wedge_x$	Introduction ( $\wedge I$ ) $\uparrow$	$\mathcal{Q}$
$\wedge_x$	Left elimination ( $\wedge E$ ) $\uparrow \downarrow$	$\mathcal{Q}$
$\wedge_x$	Right elimination ( $\wedge E$ ) $\uparrow \downarrow$	$\mathcal{Q}$
Disjunction		
$\vee_x$	Left introduction ( $\vee I$ ) $\uparrow$	$\mathcal{Q}$
$\vee_x$	Right introduction ( $\vee I$ ) $\uparrow$	$\mathcal{Q}$
$\vee_x$	Elimination ( $\vee E$ ) $\uparrow \downarrow$	$\mathcal{Q}$
Negation		
$\neg_x$	Introduction ( $\neg I$ ) $\uparrow$	$\mathcal{Q}$
$\neg_x$	Elimination ( $\neg E$ ) $\uparrow$	$\mathcal{Q}$
False		
$\perp_x$	Elimination ( $\perp E$ ) $\uparrow$	$\mathcal{Q}$

(1)	$P \vee (Q \wedge R)$	<i>hypothesis</i>
(2)	$P$	<i>hypothesis</i>
(3)	$P \vee Q$	<b>to be justified</b>
(4)	$Q \wedge R$	<i>hypothesis</i>
5	$P \vee Q$	<b>to be justified</b>
(6)	$P \vee Q$	(1) (2) ... (3) (4) ... (5) <i>Admitted</i>
(7)	$(P \vee (Q \wedge R)) \Rightarrow (P \vee Q)$	(1) ... (6) ( $\Rightarrow I$ )

# Edukera (math mode)

Home Analysis Induction Exercise 1

Let  $P$  be a proposition defined at rank  $n$  by  $\sum_{k=0}^n k = \frac{n \cdot (n+1)}{2}$  definition

(1)  $P(0)$  to be justified

Let  $n$  be a natural integer declaration

(2)  $P(n)$

3  $\sum_{k=0}^n k = \frac{n \cdot (n+1)}{2}$

(4)  $\sum_{k=0}^{n+1} k = \frac{(n+1) \cdot (n+2)}{2}$

(5)  $P(n+1)$

(6) For every natural integer  $n$ ,  $P(n)$

(7) For every natural integer  $n$ ,  $\sum_{k=0}^n k = \frac{n \cdot (n+1)}{2}$  (6) by definition of  $P$

Deduction from (3)

1 2

(3)  $\sum_{k=0}^n k = \frac{n \cdot (n+1)}{2}$  reference

a  $\left( \sum_{k=0}^n k \right) + 1? = \frac{n \cdot (n+1)}{2} + 1?$  to (3) by adding 1? to both sides

To be justified: (1) (4)

Value of 1? :  $n+1$

P	n	7	8	9	/	$\pi$	$+\infty$	$+$	max	min	$\wedge$	$x^2$	U	$\emptyset$	R	$R^*$
		4	5	6	$\times$		$-\infty$	$-$	ln	exp	$\sqrt{\quad}$	$x^{-1}$	n	u	$R^+$	$R^-$

# Edukera (prototype for geometry)

Home | Calculus | Exercise 12 | edukera

**Deduction of** [X]

Prove: Reasoning  
parallelism of opposite sides  
non-contradiction  
equal opposite sides  
equality of opposite vectors

Initial content: 1 2

Let A  
Let B Preview  
Let C  
Let I Conclusion  
Let J  
1 Assum according to 5, by parallelism of opposite sides  
2 Assum [Apply]

Let K be a point  
3 K is the symmetric of J with respect to I by construction of K  
4 I is the midpoint of the segment [ KJ ] according to 3, by definition of the symmetric  
5 AKBJ is a parallelogram [Done] according to 4 1, by intersection of the diagonals at their mutual midpoint I  
Conclusion  
6 ( BC ) and ( IJ ) are parallel

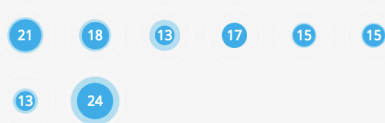
E 12  
onstruction : (section 1.3, auto translated) into practice in elementary geometry.

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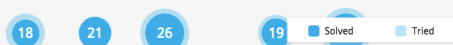
Classical logic



Distributive properties



De Morgan's laws

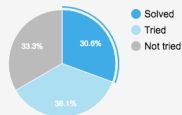


Exercise 29

Prove that

for every propositions A B C,

$$((A \Rightarrow (B \Rightarrow C)) \Rightarrow C) \Rightarrow (((A \Rightarrow C) \Rightarrow C) \Rightarrow ((B \Rightarrow C) \Rightarrow C) \Rightarrow C)$$



Étudiants (11)

	First name	Last name	Temps
	...	...	41:27
	...	...	1:13:54
	...	...	16:29

temps moyen: 2:55:20 Progression 78%

# Some experiments

Previous years:

- Undergraduate computer-science logic course: natural deduction (Edukera/Logic Mode)
- Graduate computer-science formal theorem proving course (Edukera Logic Mode+Coq)
- Graduate computer-science software-foundations course (Coq, Frama-c, why3)

Starting September:

- First-year undergraduate maths/computer science: The concept of proof and very basics results about relations/functions/sets (Edukera Maths Mode).

# Results in a logic course

- 36 students, > 2000 exercises in natural deduction
- positive student feedback
- need a scientific evaluation



# Results in a course about formal theorem proving

- Edukera as a tool to learn natural deduction.
- Coq tactics are then learned quicker.

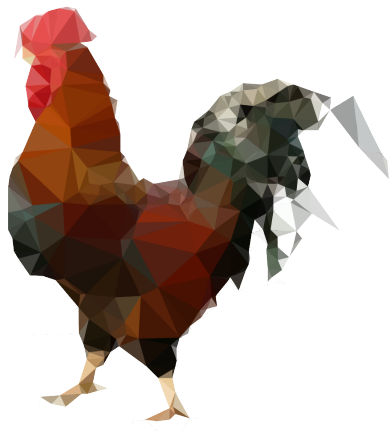
# Outline

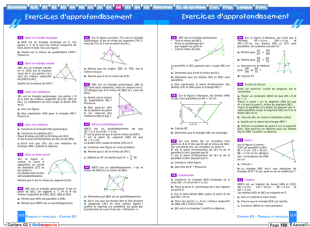
## 1 Interactive Theorem Proving for the Education

## 2 Overview of GeoCoq

- Foundations
- Arithmetization of Geometry
- Automation
- Continuity
- Logic
- 34 parallel postulates
- Two formalizations of the Elements
- Some high-school examples

- An Open Source library about foundations of geometry
- Michael Beeson, Gabriel Braun, Pierre Boutry, Charly Gries, Julien Narboux, Pascal Schreck
- Size: > 3900 Lemmas,  
> 130000 lines
- License: LGPL3

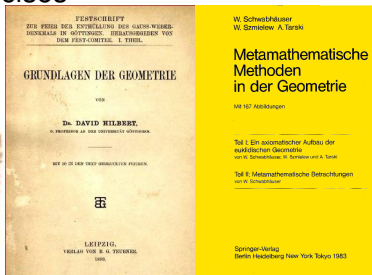




## Exercices



Euclide



Hilbert

Tarski

## What we have:

**Axiom systems** Tarski's, Hilbert's, Euclid's and variants.

**Foundations** In arbitrary dimension, in neutral geometry.  
Between-ness, Two-sides, One-side, Collinearity,  
Midpoint, Symmetric point, Perpendicularity, Parallelism,  
Angles, Co-planarity, . . .

**Classic theorems** Pappus, Pythagoras, Thales' intercept theorem,  
Thales' circle theorem, nine point circle, Euler line,  
orthocenter, circumcenter, incenter, centroid,  
quadrilaterals, Sum of angles, Varignon's theorem, . . .

**Arithmetization** Coordinates

**High-school** Some exercises

## What is missing:

- Consequence of continuity: trigonometry, areas
- link with Complex numbers

# Foundations of geometry

- 1 Synthetic geometry
- 2 Analytic geometry
- 3 Metric geometry
- 4 Transformations based approaches

# Synthetic approach

Assume some undefined geometric objects + geometric predicates + axioms ...

The name of the assumed types are not important.

- Hilbert's axioms:

types: points, lines and planes

predicates: incidence, between, congruence of segments, congruence of angles

- Tarski's axioms:

types: points

prédicats: between, congruence

- ... many variants

# Analytic approach

We assume we have numbers (a field  $\mathbb{F}$ ).

We define geometric objects by their coordinates.

Points :=  $\mathbb{F}^n$

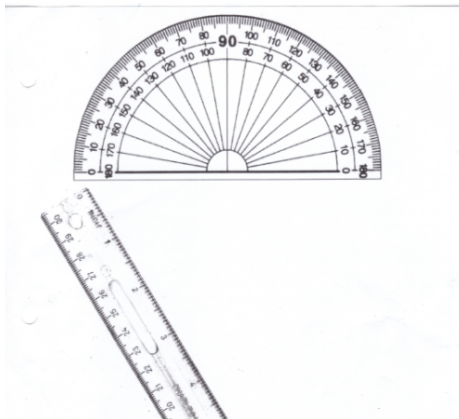


# Metric approach

Compromise between synthetic and metric approach.

We assume both:

- numbers (a field)
- geometric objects
- axioms



- Birkhoff's axioms: points, lines, reals, ruler and protractor
- Chou-Gao-Zhang's axioms: points, numbers, three geometric quantities









# Transformation groups

Erlangen program. Foundations of geometry based on group actions and invariants.

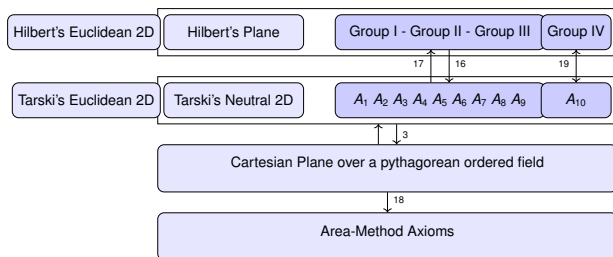


Felix Klein

# Comparison

	Synthetic	Analytic
Logical Reasoning		
Proof reuse between geometries		
Computations		
Automatic proofs		

# Overview of the axiom systems



<sup>16</sup>Gabriel Braun, Pierre Boutry, and Julien Narboux (2016). "From Hilbert to Tarski". In: [Eleventh International Workshop on Automated Deduction in Geometry](#). Proceedings of ADG 2016

<sup>17</sup>Gabriel Braun and Julien Narboux (2012). "From Tarski to Hilbert". English. In: [Post-proceedings of Automated Deduction in Geometry 2012](#). Vol. 7993. LNCS

<sup>18</sup>Pierre Boutry, Gabriel Braun, and Julien Narboux (2017). "Formalization of the Arithmetization of Euclidean Plane Geometry and Applications". In: [Journal of Symbolic Computation](#)

<sup>19</sup>**boutry'parallel'2015**

# An "axiom free" development

Axiom = global variable

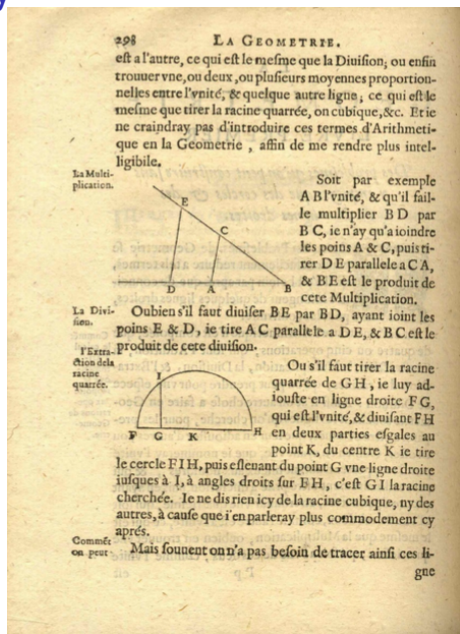
```
Class Tarski_neutral_dimensionless :=
{
  Tpoint : Type;
  Bet : Tpoint -> Tpoint -> Tpoint -> Prop;
  Cong : Tpoint -> Tpoint -> Tpoint -> Tpoint -> Prop;
  cong_pseudo_reflexivity : forall A B, Cong A B B A;
  cong_inner_transitivity : forall A B C D E F,
    Cong A B C D -> Cong A B E F -> Cong C D E F;
  cong_identity : forall A B C, Cong A B C C -> A = B;
  segment_construction : forall A B C D,
    exists E, Bet A B E /\ Cong B E C D;
  ...
}
```

Then, we can also formalize some meta-theoretical results:

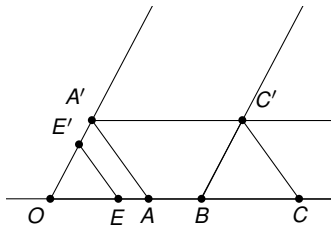
```
Instance Hilbert_euclidean_follows_from_Tarski_euclidean :  
  Hilbert_euclidean  
  Hilbert_neutral_follows_from_Tarski_neutral.
```

# Arithmetization of Geometry

René Descartes (1925).  
La géométrie.

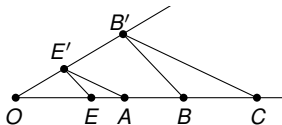


# Addition





# Multiplication



# Characterization of geometric predicates

Geometric predicate	Characterization
$AB \equiv CD$	$(x_A - x_B)^2 + (y_A - y_B)^2 - (x_C - x_D)^2 + (y_C - y_D)^2 = 0$
Bet $A B C$	$\exists t, 0 \leq t \leq 1 \wedge \begin{matrix} t(x_C - x_A) = x_B - x_A \\ t(y_C - y_A) = y_B - y_A \end{matrix} \wedge$
Col $A B C$	$(x_A - x_B)(y_B - y_C) - (y_A - y_B)(x_B - x_C) = 0$
$I$ midpoint of $AB$	$\begin{matrix} 2x_I - (x_A + x_B) = 0 \\ 2y_I - (y_A + y_B) = 0 \end{matrix} \wedge$
Per $ABC$	$(x_A - x_B)(x_B - x_C) + (y_A - y_B)(y_B - y_C) = 0$
$AB \parallel CD$	$\begin{matrix} (x_A - x_B)(x_C - x_D) + (y_A - y_B)(y_C - y_D) = 0 \\ (x_A - x_B)(x_A - x_B) + (y_A - y_B)(y_A - y_B) \neq 0 \\ (x_C - x_D)(x_C - x_D) + (y_C - y_D)(y_C - y_D) \neq 0 \end{matrix} \wedge$
$AB \perp CD$	$\begin{matrix} (x_A - x_B)(y_C - y_D) - (y_A - y_B)(x_C - x_D) = 0 \\ (x_A - x_B)(x_A - x_B) + (y_A - y_B)(y_A - y_B) \neq 0 \\ (x_C - x_D)(x_C - x_D) + (y_C - y_D)(y_C - y_D) \neq 0 \end{matrix} \wedge$

# Formalization technique: bootstrapping

**Manually** bet, cong, equality, col

**Automatically** midpoint, right triangles, parallelism and perpendicularity

# Using automation

Using Gröbner's bases, but this is not a theorem about polynomials:

```
Lemma centroid_theorem : forall A B C A1 B1 C1 G,  
  Midpoint A1 B C ->  
  Midpoint B1 A C ->  
  Midpoint C1 A B ->  
  Col A A1 G ->  
  Col B B1 G ->  
  Col C C1 G \ / Col A B C.
```

Proof.

```
intros A B C A1 B1 C1 G; convert_to_algebra; decompose_coordinates.  
intros; spliter. express_disj_as_a_single_poly; nsatz.  
Qed.
```

# Continuity properties

- Dedekind

# Continuity properties

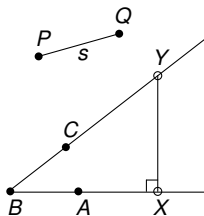
- Dedekind
- Archimedes

# Continuity properties

- Dedekind
- Archimedes
- Aristotle

# Continuity properties

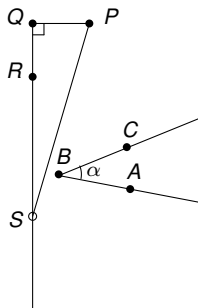
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- Dedekind
- Archimedes
- Aristotle
- Greenberg



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# Continuity properties

- Dedekind



- Archimedes



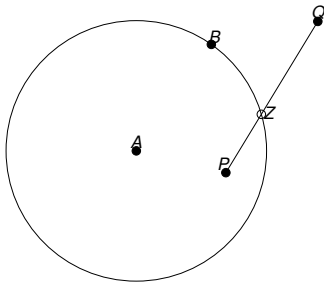
- Aristotle



- Greenberg

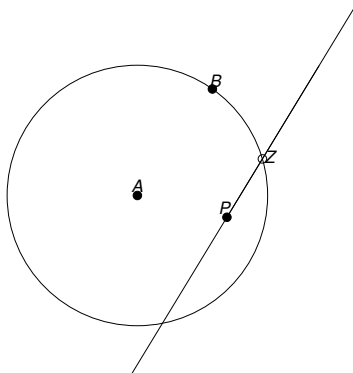
# Segment-Circle / Line-Circle continuity

- Circle-Segment



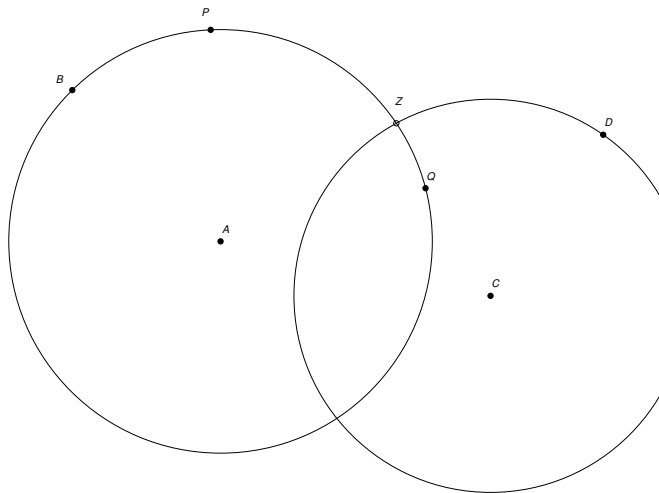
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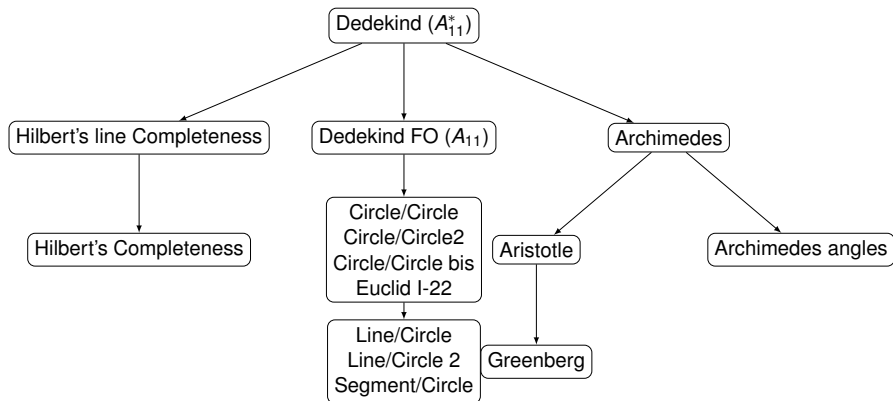
- Circle-Segment
- Circle-Line



# Segment-Circle / Line-Circle continuity

- Circle-Segment
- Circle-Line
- Circle-Circle





Continuity	Axiom
circle/line continuity	ordered Pythagorean field <sup>20</sup>
FO Dedekind cuts	ordered Euclidean field <sup>21</sup>
Dedekind	real closed field <sup>22</sup>
	reals

---

<sup>20</sup>the sum of squares is a square

<sup>21</sup>positive are square

<sup>22</sup> $F$  is euclidean and every polynomial of odd degree has at least one root in  $F$ .



## Intuitionist logic <sup>23</sup>

- Assuming :  $\forall A, B : \text{Points}, A = B \vee A \neq B$
- We prove : excluded middle for all other predicates,

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<sup>23</sup>Pierre Boutry et al. (2014). “A short note about case distinctions in Tarski’s geometry”. In:

[Proceedings of the 10th Int. Workshop on Automated Deduction in Geometry. Vol. TR 2014/01. Proceedings of ADG 2014](#)

## Intuitionist logic <sup>23</sup>

- Assuming :  $\forall A, B : \text{Points}, A = B \vee A \neq B$
- We prove : excluded middle for all other predicates, **except line intersection**

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<sup>23</sup>Pierre Boutry et al. (2014). “A short note about case distinctions in Tarski’s geometry”. In:

[Proceedings of the 10th Int. Workshop on Automated Deduction in Geometry. Vol. TR 2014/01. Proceedings of ADG 2014](#)

# Outline

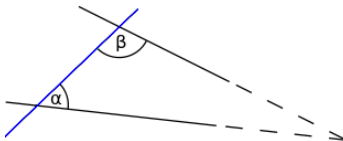
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# Euclid 5th postulate

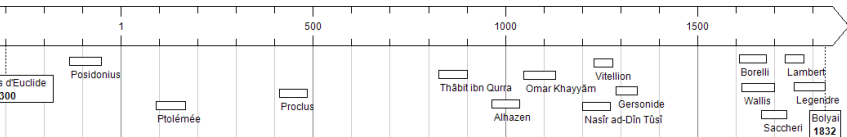
*“If two lines are drawn which intersect a third in such a way that the sum of the inner angles on one side is less than two right angles, then the two lines inevitably must intersect each other on that side if extended far enough.”*



# History

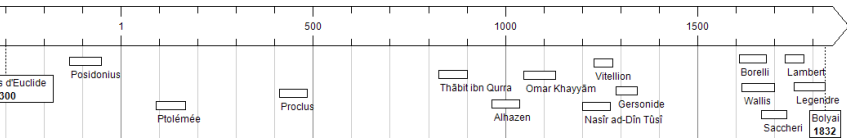
- A less obvious postulate

# History

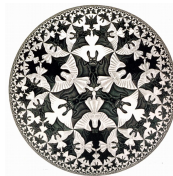


- A less obvious postulate
- Incorrect proofs during centuries

# History

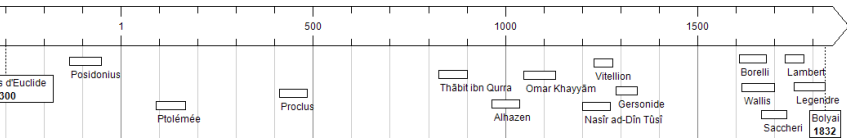


- A less obvious postulate
- Incorrect proofs during centuries
- Independence

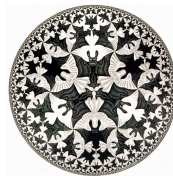


Escher, Circle Limit IV, 1960

# History



- A less obvious postulate
- Incorrect proofs during centuries
- Independence
- Some equivalent statements



Escher, Circle Limit IV, 1960



# A long history of incorrect proofs . . .

In 1763, Klügel <sup>24</sup> provides a list of 30 failed attempts at proving the parallel postulate.

## Examples:

- Ptolémée uses implicitly Playfair's postulate (uniqueness of the parallel).
- Proclus uses implicitly "Given two parallel lines, if a line intersect one of them it intersects the other".
- Legendre published several incorrect proofs in its *best-seller* "Éléments de géométrie".

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<sup>24</sup> [G. S. Klugel \(1763\)](#). "Conatum praecipuorum theoriam parallelarum demonstrandi recensio". [PhD thesis. Schultz, Göttingen](#)

# Mistakes

- Circular arguments

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▶ `parallelogram ABCD := AB || CD ∧ AD || BC`

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  - ▶  $\text{parallelogram } ABCD := AB \parallel CD \wedge AD \parallel BC$
  - ▶  $\text{parallelogram2 } ABCD := AB \parallel CD \wedge AB \equiv CD \wedge$   
 $\text{Convex } ABCD$

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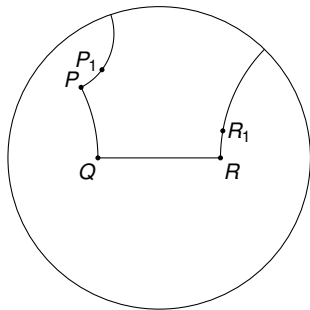
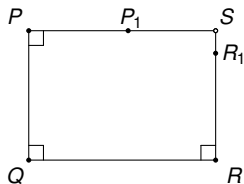
## Warning !

$(\text{parallelogram2 } ABCD \Leftrightarrow \text{parallelogram2 } BCDA) \Leftrightarrow$   
 $\text{Euclid5}$

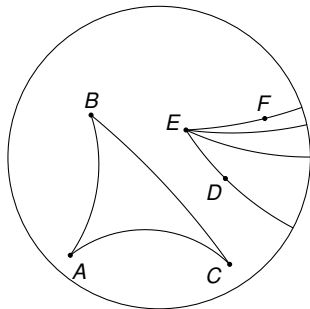
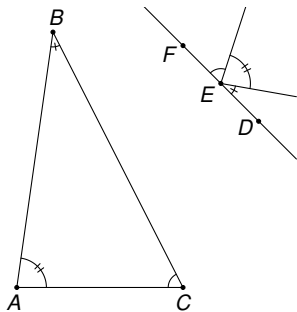


# Bachmann's Lotschnittaxiom

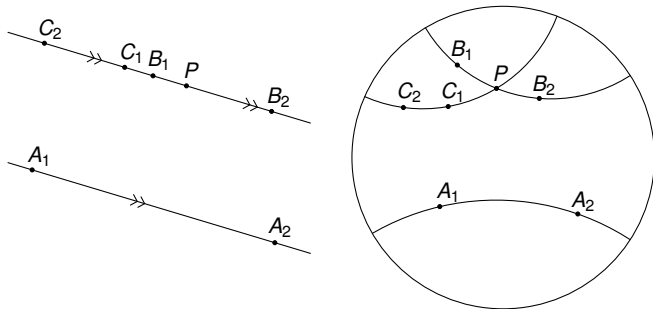
If  $p \perp q$ ,  $q \perp r$  and  $r \perp s$  then  $p$  and  $s$  meet.



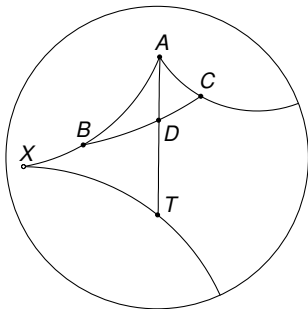
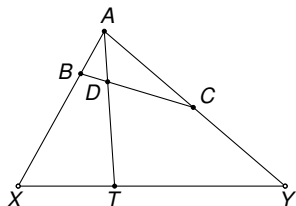
# Triangle postulate



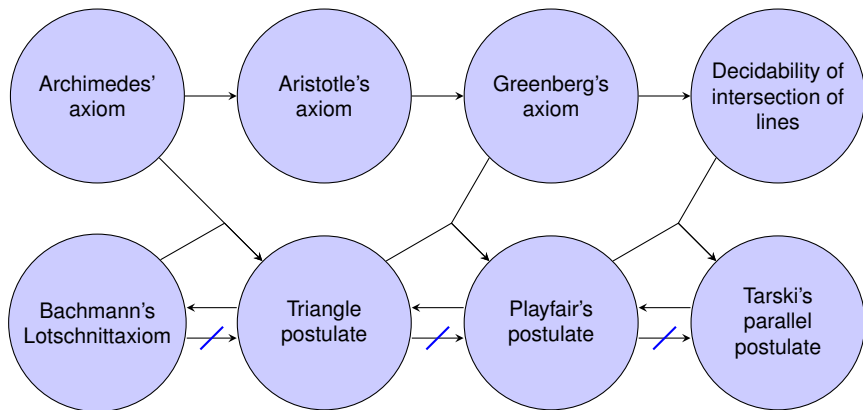
# Playfair's postulate



# Tarski's postulate



# Four groups





# Outline

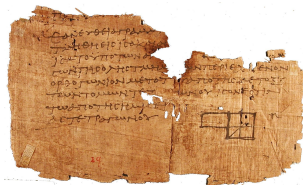
## 1 Interactive Theorem Proving for the Education

## 2 Overview of GeoCoq

- Foundations
- Arithmetization of Geometry
- Automation
- Continuity
- Logic
- 34 parallel postulates
- Two formalizations of the Elements
- Some high-school examples

# The Elements

- A very influential mathematical book (more than 1000 editions).
- First known example of an axiomatic approach.



Book 2, Prop V, Papyrus  
d'Oxyrhynchus (year 100)



Euclid



# First project

- Joint work with Charly Gries and Gabriel Braun
- Mechanizing proofs of Euclid's **statements**
- Not Euclid's proofs!
- Trying to minimize the assumptions:
  - ▶ Parallel postulate
  - ▶ Elementary continuity
  - ▶ Archimedes' axiom

# Second project

- Joint work with Michael Beeson and Freek Wiedijk <sup>25</sup>
- Formalizing Euclid's **proofs**
- A not minimal axiom system
- Filling the gaps in Euclid

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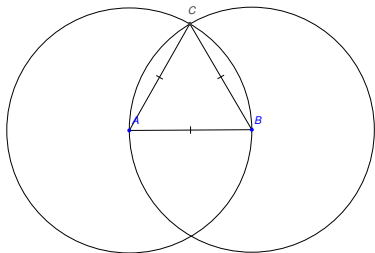
<sup>25</sup>Michael Beeson, Julien Narboux, and Freek Wiedijk (2017). “Proof-checking Euclid”.

# Example

## Proposition (Book I, Prop 1)

*Let  $A$  and  $B$  be two points, build an equilateral triangle on the base  $AB$ .*

Proof: Let  $\mathcal{C}_1$  and  $\mathcal{C}_2$  the circles of center  $A$  and  $B$  and radius  $AB$ . Take  $C$  at the intersection of  $\mathcal{C}_1$  and  $\mathcal{C}_2$ . The distance  $AB$  is congruent to  $AC$ , and  $AB$  is congruent to  $BC$ . Hence,  $ABC$  is an equilateral triangle.



# Book I, Prop 1

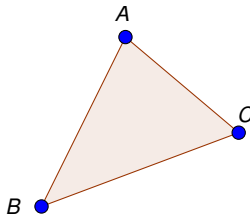
We prove two statements:

- 1 Assuming no continuity, but the parallel postulate.
- 2 Assuming circle/circle continuity, but not the parallel postulate.

Pambuccian has shown that these assumptions are minimal.

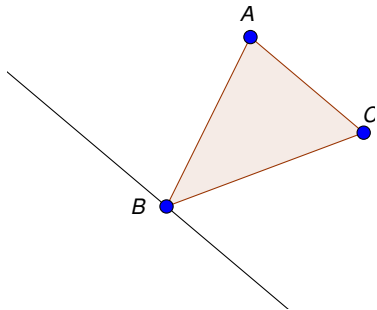
# The sum of angles of a triangle (Euclid Book I, Prop 32)

Let  $l$  be a parallel to  $AC$  through  $B$ .



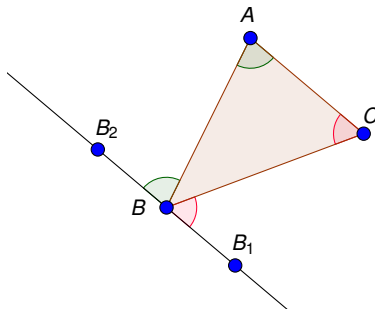
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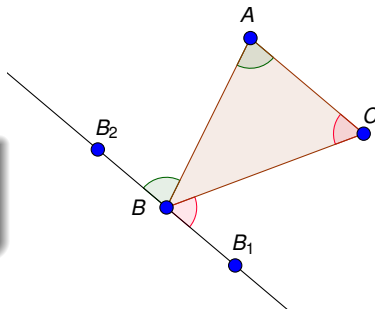


# The sum of angles of a triangle (Euclid Book I, Prop 32)

Let  $l$  be a parallel to  $AC$  through  $B$ .

**But !**

We have to prove that the angles are alternate angles.



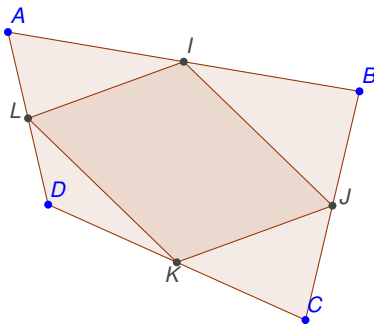


# Varignon's theorem

## Theorem

Let  $ABCD$  be a quadrilateral. Let  $I$ ,  $J$ ,  $K$  and  $L$  the midpoints of  $AB$ ,  $BC$ ,  $CD$ , and  $AD$ , then  $IJKL$  is a parallelogram.

Using the triangle midpoints theorem, in the triangle  $ABC$  we have  $AC \parallel IJ$ . We also have  $AC \parallel LK$ . Hence  $LK \parallel IJ$ . Similarly,  $IL \parallel JK$ .

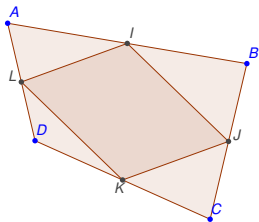


Si les côtés  $AB, BC, CD, DA$  d'une figure rectiligne de quatre côtés, sont divisés chacun en deux parties égales en  $F, G, H, E$ , & que les points des divisions soient joints par les lignes droites  $FE, EH, HG, GF$ , la figure quadrilatérale  $FEHG$  est un parallélogramme; car en menant les lignes  $DI, AC$ , comme par l'hypothèse,  $AF = FB$  &  $AI = ID$ ,  $AF : FB :: AI : ID$ , & ainsi ( *Prop. 2.* )  $EF$  est parallèle à  $DB$ . De même puisque, par l'hypothèse,  $BC = GC$ , &  $DH = HC$ ,  $BC : GC :: DH : HC$ , & par conséquent ( *Prop. 2.* )  $GH$  fera encore parallèle à la ligne  $BD$ . Donc  $EF$  &  $GH$  sont parallèles à la même troisième ligne, elles sont donc aussi parallèles entre elles.  
On peut par la même raison prouver que les lignes  $FG$  &  $EH$  sont parallèles à la ligne droite  $AC$ , & par conséquent parallèles entre elles. Donc le quadrilatère  $FEHG$  est un parallélogramme.

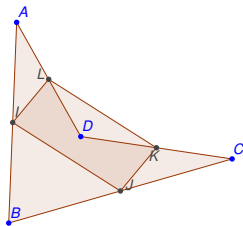
## Original proof



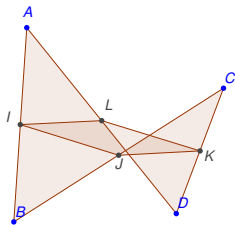
# Varignon's theorem



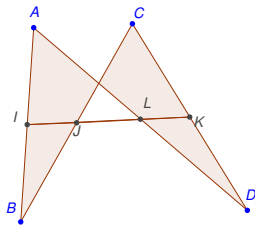
(a) Convex case



(b) Concave case



(c) Self-intersection



(d) Special case

# Challenges

- Ndgs can be easily overlooked.
- As in the Elements, text-books tend to prove properties **assuming** points in general position, but **do not check** that the points are in general position when using the properties.
- As in the Elements, text-books tend to read co-exact properties on the figure.

# Conclusion

- GeoCoq: a library for the foundations of geometry.
- Most results for high-school geometry are formalized.
- Needs integration into a GUI.
- Some challenges for automation.

Thank you